# SQL assignment DA SEPT 24 batch SUMIT GARG

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1. Create a table called employees with the following structure?

: emp\_id (integer, should not be NULL and should be a primary key) Q

: emp\_name (text, should not be NULL) Q

: age (integer, should have a check constraint to ensure the age is at least 18) Q

: email (text, should be unique for each employee) Q

: salary (decimal, with a default value of 30,000). Write the SQL query to create the above table with all constraints.

## create table employees (emp\_id int not null, emp\_name text not null, age int check(age >=18), email int unique ,

**salary dec default(30000));**

1. Explain the purpose of constraints and how they help maintain data integrity in a database. Provide examples of common types of constraints.

Constraints in a database are rules enforced on the data in tables to maintain accuracy, consistency, and integrity. They ensure that only valid and meaningful data is stored, reducing the risk of data corruption or errors.

## Purpose of Constraints

* 1. **Data Integrity**: Constraints ensure the accuracy and reliability of data.
     + Example: A constraint can ensure a person's age is always a non-negative number.
  2. **Consistency**: They enforce uniformity in data entry and relationships between tables.
     + Example: Ensuring that a product ID entered in a sales table exists in the products table.
  3. **Error Prevention**: They prevent invalid data from being entered into the database.
     + Example: Rejecting entries where a required field is missing.
  4. **Business Rules Enforcement**: Constraints can enforce specific rules relevant to an organization.
     + Example: A salary must be greater than a minimum wage.

## Common Types of Constraints

1. **Primary Key Constraint**
   * **Purpose**: Ensures each row in a table has a unique identifier.
   * **Example**: A StudentID column in a Students table must have unique values and cannot be NULL.
   * **SQL Syntax**: CREATE TABLE Students (

StudentID INT PRIMARY KEY, Name VARCHAR(50)

);

## Foreign Key Constraint

* + **Purpose**: Maintains referential integrity by ensuring that a value in one table corresponds to a value in another.
  + **Example**: A StudentID in a Enrollments table must exist in the Students table.

## SQL Syntax:

CREATE TABLE Enrollments ( EnrollmentID INT PRIMARY KEY, StudentID INT,

FOREIGN KEY (StudentID) REFERENCES Students(StudentID)

);

## Unique Constraint

* + **Purpose**: Ensures that all values in a column are unique, allowing NULL values.
  + **Example**: An email column in a Users table must have unique values.
  + **SQL Syntax**: CREATE TABLE Users (

UserID INT PRIMARY KEY, Email VARCHAR(100) UNIQUE

);

## Not Null Constraint

* + **Purpose**: Ensures that a column cannot have NULL values.
  + **Example**: A Name column in a Employees table must always have a value.
  + **SQL Syntax**: CREATE TABLE Employees (

EmployeeID INT PRIMARY KEY, Name VARCHAR(50) NOT NULL

);

## Check Constraint

* + **Purpose**: Ensures that all values in a column satisfy a specific condition.
  + **Example**: A Salary column in a Jobs table must be greater than 0.
  + **SQL Syntax**: CREATE TABLE Jobs (

JobID INT PRIMARY KEY, Salary DECIMAL(10, 2),

CHECK (Salary > 0)

);

## Default Constraint

* + **Purpose**: Assigns a default value to a column when no value is provided.
  + **Example**: A Status column in an Orders table defaults to 'Pending'.
  + **SQL Syntax**: CREATE TABLE Orders (

OrderID INT PRIMARY KEY,

Status VARCHAR(20) DEFAULT 'Pending');

3.Why would you apply the NOT NULL constraint to a column? Can a primary key contain NULL values? Justify your answer.

## Why Apply the NOT NULL Constraint to a Column?

The **NOT NULL** constraint is applied to ensure that a column always contains a value and cannot be left empty. This is crucial when the presence of data in that column is essential for the database to function correctly or to maintain data integrity.

## Reasons to Use NOT NULL:

1. **Data Completeness**: Ensures critical information is always recorded.
   * Example: A Name field in a Users table must not be NULL to identify users.
2. **Avoid Errors in Queries**: Prevents issues when performing operations or calculations that assume a value is present.
   * Example: Calculating total = quantity \* price would fail if price is NULL.
3. **Consistency in Data Relationships**: Supports referential integrity, particularly in foreign key relationships.
   * Example: A foreign key referencing a parent table cannot be NULL unless explicitly designed for optional relationships.

**Can a Primary Key Contain NULL Values?**

**No, a primary key cannot contain NULL values.**

**Justification:**

1. **Uniqueness Requirement**: The primary key uniquely identifies each row in a table. A NULL value represents "unknown" or "missing" data, which cannot serve as a unique identifier.
   * If NULL were allowed, multiple rows could have NULL values, violating the uniqueness rule.
2. **Logical Integrity**: The primary key is often used to reference rows in other tables via foreign keys. If a primary key contained NULL, the referenced rows would become ambiguous or invalid.
3. **SQL Standard**: According to SQL standards, the primary key constraint implicitly enforces both **NOT NULL** and **UNIQUE** constraints.

## Example:

The following would result in an error: CREATE TABLE Students (

StudentID INT PRIMARY KEY, Name VARCHAR(50)

);

-- Attempt to insert a NULL value into the primary key column

INSERT INTO Students (StudentID, Name) VALUES (NULL, 'John Doe'); -- Error: Primary key cannot be NULL

In conclusion:

* Apply the **NOT NULL** constraint to ensure essential data is always provided.
* A **primary key** cannot contain NULL values because it must uniquely and unambiguously identify each record in the table.

1. Explain the steps and SQL commands used to add or remove constraints on an existing table. Provide an example for both adding and removing a constraint.

## Steps and SQL Commands for Adding or Removing Constraints on an Existing Table Adding Constraints

1. **Identify the Table and Column**: Determine the table and column(s) where you want to add the constraint.
2. **Use the ALTER TABLE Command**: Use the ALTER TABLE statement with the appropriate syntax to add the constraint.
3. **Specify the Constraint Type**: Define the type of constraint (NOT NULL, UNIQUE, CHECK, etc.) and the rules it enforces.

## Example: Adding a Constraint

Suppose we have a table Employees without any constraints on the Email column. CREATE TABLE Employees (

EmployeeID INT PRIMARY KEY, Name VARCHAR(50),

Email VARCHAR(100));

To add a **UNIQUE** constraint on the Email column:

ALTER TABLE Employees

ADD CONSTRAINT UniqueEmail UNIQUE (Email);

To add a **CHECK** constraint to ensure EmployeeID is greater than 0: ALTER TABLE Employees

ADD CONSTRAINT CheckEmployeeID CHECK (EmployeeID > 0);

## Removing Constraints

1. **Identify the Constraint Name**: Determine the name of the constraint you want to remove. You can find constraint names by querying the database metadata (e.g., INFORMATION\_SCHEMA in SQL Server).
2. **Use the ALTER TABLE Command**: Use the ALTER TABLE statement with the DROP CONSTRAINT clause.

## Example: Removing a Constraint

To remove the **UNIQUE** constraint on the Email column:

ALTER TABLE Employees

DROP CONSTRAINT UniqueEmail;

To remove the **CHECK** constraint on EmployeeID:

ALTER TABLE Employees

DROP CONSTRAINT CheckEmployeeID;

1. Explain the consequences of attempting to insert, update, or delete data in a way that violates constraints. Provide an example of an error message that might occur when violating a constraint.

## Consequences of Violating Constraints

When you attempt to insert, update, or delete data in a way that violates constraints in a database, the database management system (DBMS) enforces the rules defined by the constraints and prevents the operation. This ensures data integrity but results in an error.

## Common Scenarios and Their Consequences

1. **Violation of a NOT NULL Constraint**
   * **Attempt**: Inserting or updating a row with a NULL value in a column defined as NOT NULL.
   * **Consequence**: The operation fails because the column requires a value.
   * **Example**: If the Name column in an Employees table has a NOT NULL constraint:

INSERT INTO Employees (EmployeeID, Name) VALUES (1, NULL);

**Error Message** (MySQL):

ERROR 1048 (23000): Column 'Name' cannot be null

## Violation of a UNIQUE Constraint

* + **Attempt**: Inserting or updating a row with a duplicate value in a column defined as UNIQUE.
  + **Consequence**: The operation fails because duplicate values are not allowed.
  + **Example**: If the Email column in an Users table has a UNIQUE constraint:

INSERT INTO Users (UserID, Email) VALUES (1, ['example@example.com');](mailto:%27example@example.com) INSERT INTO Users (UserID, Email) VALUES (2, ['example@example.com');](mailto:%27example@example.com) **Error Message** (PostgreSQL):

ERROR: duplicate key value violates unique constraint "users\_email\_key" DETAIL: Key [(email)=(example@example.com)](mailto:(email)%3D(example@example.com) already exists.

## Violation of a PRIMARY KEY Constraint

* + **Attempt**: Inserting or updating a row with a duplicate or NULL value in a primary key column.
  + **Consequence**: The operation fails because a primary key must be unique and cannot be NULL.

## Example:

INSERT INTO Employees (EmployeeID, Name) VALUES (1, 'Alice'); INSERT INTO Employees (EmployeeID, Name) VALUES (1, 'Bob'); **Error Message** (SQL Server):

Msg 2627, Level 14, State 1, Line 1

Violation of PRIMARY KEY constraint 'PK\_Employees'. Cannot insert duplicate key in object 'dbo.Employees'.

## Violation of a FOREIGN KEY Constraint

* + **Attempt**: Inserting or updating a row with a value in a foreign key column that does not exist in the referenced table, or deleting a referenced row in the parent table.
  + **Consequence**: The operation fails to preserve referential integrity.

## Example:

INSERT INTO Orders (OrderID, CustomerID) VALUES (1, 999); -- Assuming 999 does not exist in the Customers table

**Error Message** (MySQL):

ERROR 1452 (23000): Cannot add or update a child row: a foreign key constraint fails (`Orders`, CONSTRAINT `fk\_customer` FOREIGN KEY (`CustomerID`) REFERENCES

`Customers` (`CustomerID`))

## Violation of a CHECK Constraint

* + **Attempt**: Inserting or updating a row with a value that does not satisfy the condition defined in a CHECK constraint.
  + **Consequence**: The operation fails because the value does not meet the defined rule.
  + **Example**: If a Salary column has a CHECK constraint ensuring it must be greater than 0:

INSERT INTO Jobs (JobID, Salary) VALUES (1, -500);

**Error Message** (Oracle):

ORA-02290: check constraint (SCHEMA.CHECK\_SALARY) violated

1. You created a products table without constraints as follows: CREATE TABLE products (

product\_id INT, product\_name VARCHAR(50), price DECIMAL(10, 2));

Now, you realise that?

: The product\_id should be a primary key Q

: The price should have a default value of 50.00

ALTER TABLE products

ADD CONSTRAINT pk\_product\_id PRIMARY KEY (product\_id);

ALTER TABLE products

ALTER COLUMN price SET DEFAULT 50.00;

1. You have two tables: SELECT

students.student\_name, classes.class\_name

FROM

students

INNER JOIN

classes ON

students.class\_id = classes.class\_id;

1. Consider the following three tables:

SELECT

products.order\_id, customers.customer\_name, products.product\_name

FROM

orders LEFT JOIN

products ON

orders.order\_id = products.order\_id INNER JOIN

customers ON

orders.customer\_id = customers.customer\_id;

1. Given the following tables: SELECT

products.product\_id, products.product\_name, SUM(sales.amount) AS sales\_amt

FROM

sales

INNER JOIN

products ON sales.product\_id = products.product\_id GROUP BY products.product\_id , products.product\_name;

1. You are given three tables: SELECT

orders.customer\_id, customers.customer\_name, SUM(order\_details.quantity)

FROM

orders

INNER JOIN

order\_details ON orders.order\_id = order\_details.order\_id INNER JOIN

customers ON customers.customer\_id = orders.customer\_id GROUP BY orders.customer\_id , customers.customer\_name;

# Normalization & CTE

1. First Normal Form (1NF): a. Identify a table in the Sakila database that violates 1NF. Explain how you would normalize it to achieve 1NF.

In the Sakila database, a table that could potentially violate **First Normal Form (1NF)** is the address table. Specifically, consider a hypothetical case where the address table might have a column named phone\_numbers that stores multiple phone numbers as a comma-separated list.

This violates 1NF because:

* + **1NF Rule**: All columns must contain atomic (indivisible) values. Storing multiple values in a single column (e.g., a list of phone numbers) violates this rule.

To normalize it ; CREATE TABLE phone (

phone\_id INT AUTO\_INCREMENT PRIMARY KEY,

address\_id INT NOT NULL, phone\_number VARCHAR(15) NOT NULL,

FOREIGN KEY (address\_id) REFERENCES address(address\_id));

1. Second Normal Form (2NF): a. Choose a table in Sakila and describe how you would determine whether it is in 2NF. If it violates 2NF, explain the steps to normalize it.

**Understanding 2NF Requirements**

A table is in **2NF** if:

1. It is already in **First Normal Form (1NF)** (i.e., it has atomic values and no repeating groups).
2. Every non-prime attribute (non-key column) is fully functionally dependent on the entire primary key and not just part of it.

**Composite Key Assumption**

If the table had a composite primary key such as (inventory\_id, rental\_date), this table would need to be checked for 2NF compliance.

**Violations of 2NF**

If a non-prime attribute (e.g., staff\_id) depends only on a part of the composite primary key (inventory\_id) and not the entire composite key (inventory\_id, rental\_date), it violates 2NF.

**Example of Partial Dependency:**

* + staff\_id might be associated with inventory\_id because staff manage specific inventories.
  + This means staff\_id is not dependent on the entire primary key (inventory\_id, rental\_date), violating 2NF.

**Steps to Normalize to 2NF**

1. **Identify the Partial Dependency**:
   * Determine which non-prime attributes depend only on part of the composite key. In this case, staff\_id depends on inventory\_id.
2. **Decompose the Table**:
   * Create a separate table to store the relationship between inventory\_id and staff\_id.
3. **Revised Tables**:
   * **rental Table**:

| rental\_id | inventory\_id | customer\_id | rental\_date | return\_date |

| | | | | |

| 1 | 100 | 1 | 2023-01-01 10:00:00 | 2023-01-03 15:00:00 |

| 2 | 101 | 1 | 2023-01-02 12:00:00 | 2023-01-04 18:00:00 |

* + **inventory\_staff Table**:

| inventory\_id | staff\_id |

| | |

| 100 | 2 |

| 101 | 3 |

**SQL to Normalize**

1. **Create the inventory\_staff Table**: CREATE TABLE inventory\_staff (

inventory\_id INT NOT NULL, staff\_id INT NOT NULL, PRIMARY KEY (inventory\_id),

FOREIGN KEY (inventory\_id) REFERENCES inventory(inventory\_id), FOREIGN KEY (staff\_id) REFERENCES staff(staff\_id));

1. Third Normal Form (3NF): a. Identify a table in Sakila that violates 3NF. Describe the transitive dependencies present and outline the steps to normalize the table to 3NF.

**Third Normal Form (3NF)**

A table is in **Third Normal Form (3NF)** if:

1. It is in **Second Normal Form (2NF)**.
2. It has no **transitive dependencies**, meaning that non-prime attributes (non-key attributes) are not dependent on other non-prime attributes.

**Identifying Transitive Dependencies**

In this table:

* + **Primary Key**: film\_id
  + **Non-prime Attributes**: title, description, release\_year, language\_id, language\_name, rental\_duration, rental\_rate

There is a **transitive dependency**:

* + language\_name depends on language\_id, not directly on the primary key film\_id.
  + language\_id determines language\_name, which causes a transitive dependency: film\_id -> language\_id -> language\_name.

This violates **3NF** because language\_name is not directly dependent on film\_id but depends on language\_id (a non-prime attribute).

**Steps to Normalize to 3NF**

To normalize the film table to **3NF**, we need to remove the transitive dependency by creating a new table to store the language\_id and language\_name.

1. **Create a New language Table**:
   * The language table will store language\_id and language\_name, ensuring that language\_name is only dependent on language\_id, not on film\_id.

CREATE TABLE language ( language\_id INT PRIMARY KEY,

language\_name VARCHAR(255) NOT NULL

);

1. **Modify the film Table**:
   * Remove language\_name from the film table since it is now stored in the language table. The film table will still store language\_id as a foreign key.

CREATE TABLE film (

film\_id INT PRIMARY KEY,

title VARCHAR(255) NOT NULL,

description TEXT, release\_year YEAR, language\_id INT, rental\_duration INT, rental\_rate DECIMAL(5,2),

FOREIGN KEY (language\_id) REFERENCES language(language\_id)

);

1. **Populate the language Table**:
   * Insert distinct language\_id and language\_name combinations from the film table into the language table.

INSERT INTO language (language\_id, language\_name) SELECT DISTINCT language\_id, language\_name FROM film;

1. **Update the film Table**:
   * Remove the language\_name column from the film table. ALTER TABLE film DROP COLUMN language\_name;

**Normalized Tables in 3NF**

After normalization, the two tables are:

1. **language Table**: plaintext

Copy code

| language\_id | language\_name |

| | |

| 1 | English |

| 2 | Italian |

1. **film Table**: plaintext

Copy code

| film\_id | title | description | release\_year | language\_id | rental\_duration | rental\_rate |

| | | | | | | |

| 1 | The Matrix | Sci-fi thriller | 1999 | 1 | 7 | 4.99 |

| 2 | The Godfather | Crime drama | 1972 | 2 | 5 | 3.99 |

**Outcome**

* The **transitive dependency** (film\_id -> language\_id -> language\_name) has been removed.
* The film table now only contains attributes that are directly dependent on the primary key (film\_id), and the language table handles the relationship between language\_id and language\_name.
* Both tables are now in **Third Normal Form (3NF)**.

1. Normalization Process: a. Take a specific table in Sakila and guide through the process of normalizing it from the initial unnormalized form up to at least 2NF.

Let's walk through the **normalization process** for a specific table in the Sakila database. We will choose the **rental** table for this example and guide you through the steps to normalize it from **unnormalized form** to **Second Normal Form (2NF)**.

**Starting with the Unnormalized Form (UNF)**

In the unnormalized form, data is typically stored with repeating groups and non-atomic values. For example, let's assume the rental table contains both rental information and customer details in a single row:

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| **rental\_i**  **d** | **inventory\_i**  **d** | **customer\_i**  **d** | **customer\_nam**  **e** | **customer\_addres**  **s** | **rental\_dat**  **e** | **return\_dat**  **e** | **staff\_i**  **d** |
| 1 | 100 | 1 | John Doe | 123 Elm St. | 2023-01-01  10:00 | 2023-01-03  15:00 | 2 |
| 2 | 101 | 2 | Jane Smith | 456 Oak Ave. | 2023-01-02  12:00 | 2023-01-04  18:00 | 3 |

In this unnormalized form, we have:

1. **Repeating groups**: Customer details (like customer\_name and customer\_address) are repeated for each rental.
2. **Non-atomic values**: Customer details should be stored in separate fields to ensure atomicity.

**Step 1: Convert to First Normal Form (1NF)**

To convert this table into **1NF**, we need to:

* + Ensure that all columns contain atomic values (no repeating groups or multi-valued attributes).
  + Each record should represent a unique entity, so we’ll remove repeating customer data for each rental by separating customer information into another table.

**1NF Design**

We'll split the customer information into a separate customer table and keep the rental details in the rental table.

**1NF Tables**:

1. **customer Table**:

| customer\_id | customer\_name | customer\_address |

| | | |

| 1 | John Doe | 123 Elm St. |

| 2 | Jane Smith | 456 Oak Ave. |

1. **rental Table**:

| rental\_id | inventory\_id | customer\_id | rental\_date | return\_date | staff\_id |

|  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- |
| | | | | | | | | | | | | | |  |
| | 1 | | 100 | | 1 | | 2023-01-01 10:00:00 | 2023-01-03 15:00:00 | 2 | | | | | |
| | 2 | | 101 | | 2 | | 2023-01-02 12:00:00 | 2023-01-04 18:00:00 | 3 | | | | | |

Now, the **rental** table has only atomic values, and the repeating customer data has been moved to the

**customer** table. This satisfies **1NF**.

**Step 2: Convert to Second Normal Form (2NF)**

To convert the tables into **2NF**, we need to:

* + Ensure that the table is in **1NF**.
  + Remove any **partial dependencies** (where non-prime attributes depend on only part of a composite primary key).

In the **rental** table, the primary key is rental\_id, which uniquely identifies each rental. However, in the unnormalized design, customer\_id (a non-prime attribute) is dependent on rental\_id, and customer- related information like customer\_name and customer\_address are repeated for every rental.

The **rental** table is in **1NF**, but it violates **2NF** due to the partial dependency:

* + customer\_name and customer\_address depend on customer\_id (not the entire primary key, rental\_id).

**Solution to Achieve 2NF:**

We need to separate customer-specific information that depends solely on customer\_id into a new table. The rental table should only store rental-related information, referencing customer\_id as a foreign key.

**2NF Design**:

1. **customer Table**:

| customer\_id | customer\_name | customer\_address |

| | | |

| 1 | John Doe | 123 Elm St. |

| 2 | Jane Smith | 456 Oak Ave. |

1. **rental Table**:

| rental\_id | inventory\_id | customer\_id | rental\_date | return\_date | staff\_id |

| | | | | | |

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| | 1 | | 100 | | 1 | | 2023-01-01 10:00:00 | 2023-01-03 15:00:00 | 2 | | |
| | 2 | | 101 | | 2 | | 2023-01-02 12:00:00 | 2023-01-04 18:00:00 | 3 | | |

Here, the rental table now only stores rental-related data, with customer\_id being a foreign key. All customer details are stored in the customer table. The partial dependency has been removed, and the table is now in **2NF**.

**Conclusion**

* + The original unnormalized rental table contained repeating groups and non-atomic values.
  + After applying the steps to **1NF**, we separated customer data into the customer table, removing the repeating customer details.
  + We then applied **2NF** by removing the partial dependency of customer data on rental\_id and putting customer information in its own table.

The table is now normalized to **2NF**, ensuring all non-prime attributes are fully dependent on the primary key.